# Calculating Szeged Index and Revised Szeged Index by Using Adjacency Matrix 

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#### Abstract

Topological indices are graph invariants used in theoretical chemistry to encode molecules for the design of chemical compounds with given physicochemical properties or given pharmacological and biological activities. Szeged index $(S z(G))$ and Revised Szeged index $\left(S z^{*}(G)\right)$ of molecular graph identify some distance properties for graphs. In computational chemistry and graph theory, $(S z(G))$ and $\left(S z^{*}(G)\right)$ were more available to test the characteristics of chemical molecular structures, and thus widely used in chemical applications. In this paper, a simple algorithm is presented for constructing the distance matrix. This algorithm is designed for calculation of $S z(G)$ and $S z^{*}(G)$.


## Keywords

Adjacency Matrix; Connected Graph; Distance Matrix; Revised Szeged Index; Szeged Index.

## 1. INTRODUCTION

Suppose $G=(V, E)$ be a simple and connected graph, with vertex set $V$ and edge set $E$. Let $n=$ $|V|$ and $m=|E|$ denote the order and the size of $G$.
The distance between two vertices $u$ and $v$ in graph is the number of edges in the shortest path and is denoted by $d(u, v)$.
Let $e=u v \in E$ and define the partition, with respect to $e,\left\{N_{u}(e), N_{v}(e), N_{0}(e)\right\}$ of the vertices of $G$ as follows [1]:

$$
\begin{aligned}
& N_{u}(e)=\{w \in V: d(u, w)<d(v, w)\}, \\
& N_{v}(e)=\{w \in V: d(v, w)<d(u, w)\}, \\
& N_{0}(e)=\{w \in V: d(v, w)=d(u, w)\} .
\end{aligned}
$$

Let $n_{u}(e), n_{v}(e)$ and $n_{0}(e)$ denote the number of vertices in $N_{u}(e), N_{v}(e)$ and $N_{0}(e)$, respectively. Szeged index [2,3] of $G$ is defined by
$S z=S z(G)=\sum_{e=u v \in E} n_{u}(e) . n_{v}(e)$.
The definition of the Szeged index does not take into account the vertices at equal distance to $u$ and $v$. Szeged star index (or revised Szeged index) of $G$ is defined by $[4,5]$ :

$$
\begin{aligned}
S z^{*}= & S z^{*}(G)= \\
& \sum_{e=u v \in E}\left(n_{u}(e)+\frac{n_{0}(e)}{2}\right) \cdot\left(n_{v}(e)+\frac{n_{0}(e)}{2}\right) .
\end{aligned}
$$

Some properties and applications of Szeged index and revised Szeged index have been reported in [1, $6,7,8,9,10]$.
This paper is organizede in the following way:

After introduction, the new algorithm for obtaining $S z(G)$ and $S z^{*}(G)$ is represented in Section 2. In Section 3, effectivness of the proposed approach is verified by solving numerical examples. Finally, conclusions are discussed in the last section.

## 2. STATEMENT OF THE PROBLEM

The standard distance matrix or the vertex-distance matrix (or the minimum path matrix) of a vertexlabeled connected graph $G$, denoted by $D$, is a real symmetrix $n \times n$ matrix whose elements are defined as:

$$
D(i, j)=\left\{\begin{array}{lc}
l(i, j) & \text { if } \quad i \neq j \\
0 & \text { otherwise },
\end{array}\right.
$$

where $l(i, j)$ is the length of the shortest path, i. e., the minimum number of edges, between vertices $i$ and $j$ in $G$.
This matrix has been used to generate a number of topological indices, e.g., Balaban index [11, 12], Wiener index [13, 14], multiplicative Wiener index [15, 16, 17] and distance-sum index [18]. An efficient algorithm is available for computing the vertex-distance matrix of any graph in [19].
A common way to obtain Matrix $D$ is to use different powers of adjacency Matrix $A$.
The elements of Matrix $A$ give the connections between vertices. Powers of adjacency matrix are concatenating walks. The $i j-t h$ entry of the $k-$ th power of $A$ counts the number of walks of length $k$ from vertex $i$ to vertex $j$, not paths (a walk can repeat vertices, while a path cannot). So, to
create a distance matrix, we need to iterativerly power adjacency Matrix $A$, and as soon as a $i j-$ th element is non-zero, we have to assign the distance $k$ in distance matrix [20].

```
A = input('Enetr the adjacency matrix');
\([n, n]=\operatorname{size}(A)\);
\(D=\operatorname{NaN}(n)\);
\(B=A ;\)
\(k=1\);
while any(isnan(D(:)))
(Check for new walks and assign distance)
\(D(B>0 \& \operatorname{isnan}(D))=k\);
(Iteration)
\(k=k+1\);
\(B=B * A\);
end;
\((D(i, i)=2, \quad\) in this algorithm \()\)
for \(i=1: n\)
\(D(i, i)=0\);
end
```

In all the mentioned methods, matrix multiplication is used [19, 20].
In this paper, a simple and practical method for calculating Matrix $D$, without the need for matrix multiplication will be presented. The advantage of this method is that it can be used for finite graphs with large dimensions and no matrix storage space is necessary.
Matrix $D$ is symmetric with zero diagonal elements. Therefore, in the proposed algorithm, only the elements above diagonal are calculated. Then, the matrix can be obtained by considering symmetry.

## Algorithm:

1. Enter adjacency Matrix $A$,
2. Consider

$$
D(i, j)=A(i, j)
$$

In this step, all paths with lenght of one are defined.
3. Now starting from $h=1$ and until all the values above diagonal of Matrix $D$ become non-zero
for $\quad k=1: n$
for $\quad i=k+1: n$
for $\quad j=i+1: n$
if $\quad(A(i, j)=1 \quad \& D(k, i)=h \quad \& D(k, j)=0)$

$$
\begin{gathered}
D(k, j)=h+1 ; \\
h=h+1 ;
\end{gathered}
$$

Thus, for each $h$, all paths of length $h+1$ will be specified simultaneously and if all the values above diagonal of Matrix $D$ are non-zero, the algorithm stops.

In [21] theorem 2.2, to calculate $S z^{*}(G)$, a formula has been presented based on the number of vertices, the number of edges, $N_{u}, N_{v}$ and $S z(G)$. To use this formula, we need to calculate $N_{u}, N_{v}$ and $S z(G)$. Additionally, in corollary 2.3, for values $S z(G)$ and $S z^{*}(G)$, an upper bound has been identified on the basis of the number of edges and vertices of the graph. In this paper, we will present a simple and applicable algorithm for calculating $N_{u}, N_{v}, S z(G)$ and $S z^{*}(G)$ based on distance Matrix D.
Now, with the help of Matrix $D$ and the following algorithm, we can simply calculate $S z(G)$ and $S z^{*}(G)$.
In [8], values $S z(G)$ and $S z^{*}(G)$ for some specific graphs have been achieved based on values of Wiener index. However, for calcilating Wiener index, we need to identify $d(u, v)$ (the number of edges on any of the shortest paths joining vertex $u$ to vertex $v$ ) and in this papers, we have presented the method of calculating $d(u, v)$, [8] has not presented a method for calulating it.
Algorithm for calculating $S z(G)$ and $S z^{*}(G)$ :
Let

$$
S z=0, \quad S z^{*}=0, \quad n_{i}=0, \quad n_{j}=0, \quad n_{0}=0 .
$$

For all the values above diagonal of Matrix $D$ with $D(i, j)=1$, (This means that there is an edge $e$ between vertices $i$ and $j$ ):

1. for $\mathrm{k}=1: \mathrm{n}$
calculate $D(i, k)-D(j, k)$,
2. If $D(i, k)-D(j, k)<0$,
$n_{i}=n_{i}+1$,
(Vertices whose distance from vertex $v_{j}$ is greater than their distance from vertex $v_{i}$ )
3. If $D(i, k)-D(j, k)>0$,
$n_{j}=n_{j}+1$,
4. If $\quad D(i, k)-D(j, k)=0$,
$n_{0}=n_{0}+1$,
We do not need Step 3 to calculate $S z$.

$$
\left\{\begin{array}{l}
S z^{*}=S z^{*}+\left(n_{i}+\frac{n_{0}}{2}\right) \cdot\left(n_{j}+\frac{n_{0}}{2}\right) ; \\
S z=S z+\left(n_{i}\right) \cdot\left(n_{j}\right)
\end{array}\right.
$$

## 3. NUMERICAL EXAMPLES

In this part, some examples are given to study the efficiency of the mentioned method and also to compare it with other solving methods [8].
Example 1: Table (1) shows the values of $S z(G)$ and $S z^{*}(G)$ for some graphs, respectively.
The values obtained for Complete graphs, Complete bipartite graphs and Path graphs are the same as the results found in [8].

The values obtained for Fullerene graphs are the same as those of [22]. Therefore, this algorithm is practical for calculating $S z(G)$ and $S z^{*}(G)$.

Table 1. The values of $S z(G)$ and $S z^{*}(G)$ for special

| graphs. |  |  |
| :--- | :--- | :--- |
| G | $\mathrm{Sz}(\mathrm{G})$ | $\mathrm{Sz}^{*}(\mathrm{G})$ |
| Complete graphs $\mathrm{K}_{\mathrm{n}}$ | $\frac{\mathrm{n}(\mathrm{n}-1)}{2}$ | $\frac{\mathrm{n}^{3}(\mathrm{n}-1)}{8}$ |
| $\mathrm{~K}_{7}$ | 21 | 257.25 |
| $\mathrm{~K}_{8}$ | 28 | 448.00 |
| $\mathrm{~K}_{9}$ | 36 | 729.00 |
| Complete bipartite graphs | $\mathrm{n}^{4}$ | $\mathrm{n}^{4}$ |
| $\mathrm{~K}_{\mathrm{n}, \mathrm{n}}$ |  |  |
| $\mathrm{K}_{7,7}$ | 2401 | 2401 |
| $\mathrm{~K}_{8,8}$ | 4096 | 4096 |
| $\mathrm{~K}_{9,9}$ | 6561 | 6561 |
| Path graphs $\mathrm{P}_{\mathrm{n}}$ | $\frac{\mathrm{n}^{3}-\mathrm{n}}{6}$ | $\frac{\mathrm{n}^{3}-\mathrm{n}}{6}$ |
| $\mathrm{P}_{8}$ | 84 | 84 |
| $\mathrm{P}_{9}$ | 120 | 120 |
| $\mathrm{P}_{10}$ | 165 | 165 |
| Fullerene graphs $\mathrm{C}_{\mathrm{n}}$ |  |  |
| $\mathrm{C}_{72}$ | 95412 | 139188 |
| $\mathrm{C}_{84}$ | 163896 | 222024 |
| $\mathrm{C}_{96}$ | 235236 | 331260 |

Example 2. Let $G=\operatorname{CCC}(n)$ be the crystal structure of cubic carbon. values of $S z(G)$ and $S z^{*}(G)$ of $\operatorname{CCC}(n)$ for $n=1$ and $n=2$ is equal to $S z(G)=S z^{*}(G)=192$ and $S z(G)=$ $S z^{*}(G)=4.0704 \times 10^{4}$. The values obtained for Crystal Structure Cubic Carbon are the same as those of [23].


Fig. 1. Crystal Structure Cubic Carbon CCC (1) in part (a) and CCC (2) in part (b).

## 4. CONCLUSION

Distance Matrix $D$ has been used to generate a number of topological indices. Many studies have been done to find $S z(G)$ and $S z^{*}(G)$ for different graphs. In most cases, the upper and lower bounds are specified for these values. In some graphs, such as Complete graphs and paths, these values are calculated in terms of the number of vertices. In this paper, we calculated $D$ using adjacency matrix and without multiplying the of matrix. Then, using Matrix $D$, the values of $S z(G)$ and $S z^{*}(G)$ are easily obtained. The results are completely acceptable and satisfactory to compared the mentioned references.

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## محاسبه شاخص Szeged و Szeged اصلاحشده با استفاده از ماتريس مجاورت

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